



Biparsers: Exact-Printing for Data Synchronisation

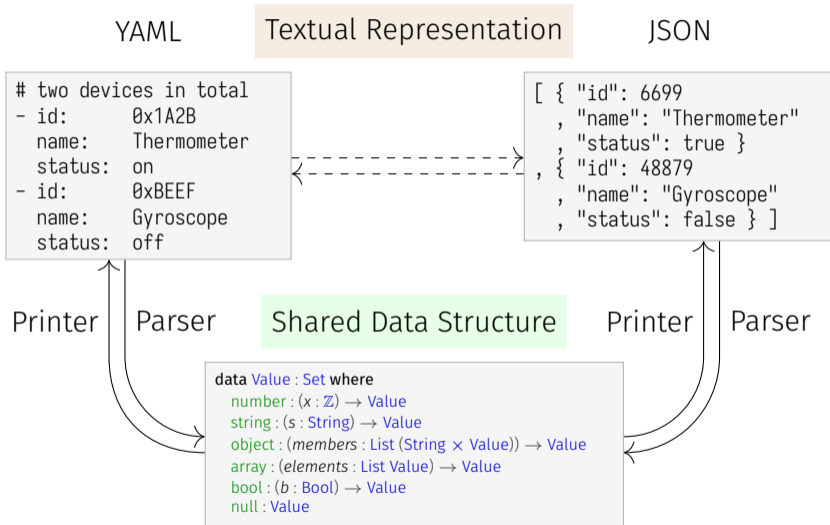
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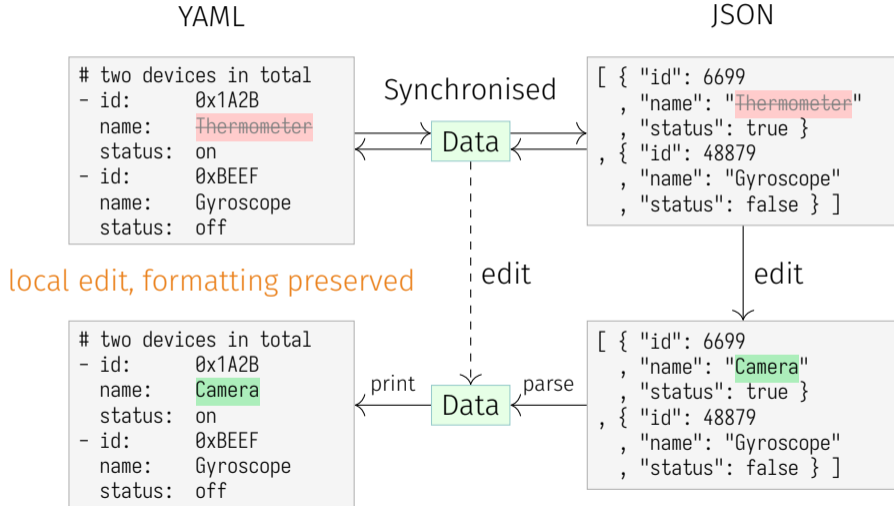
January 23, 2025 @ POPL '25

Problem: Data Synchronisation

Example: Configuration Synchronisation



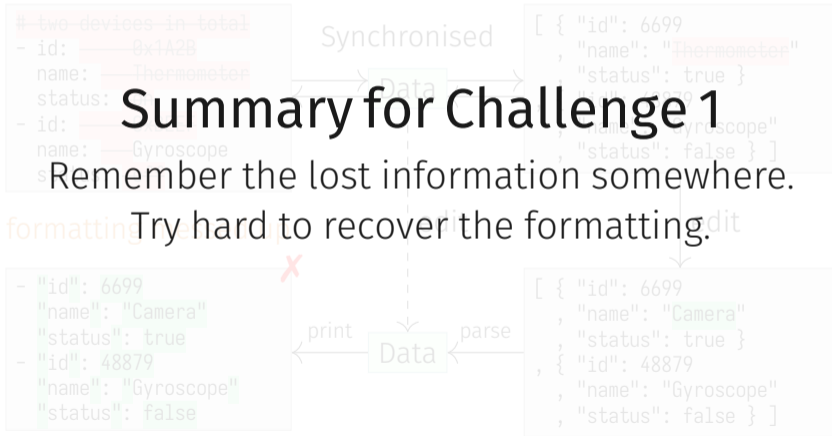
Synchronisation: Expected Behaviour



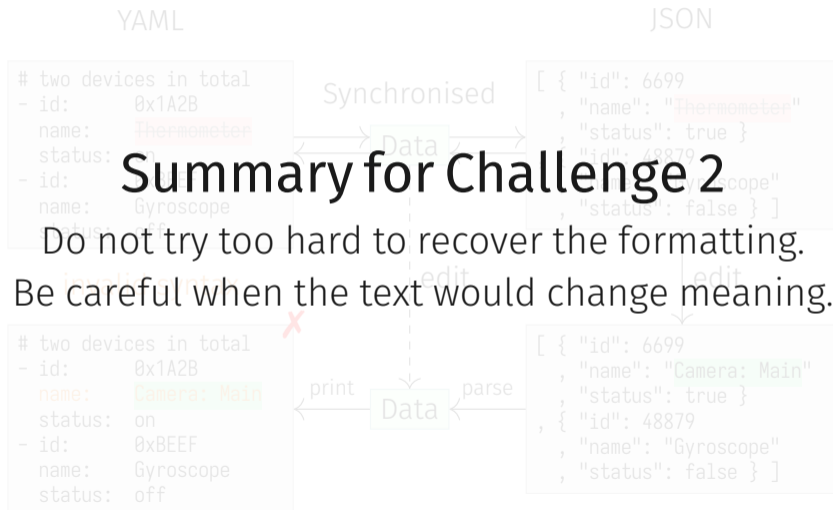
Challenge 1: Non-Injectivity (Exact-Printing)

YAML

JSON



Challenge 2: Parser-Printer Consistency



Summary for Challenge 2

Do not try too hard to recover the formatting.
Be careful when the text would change meaning.

Solution: Biparsers

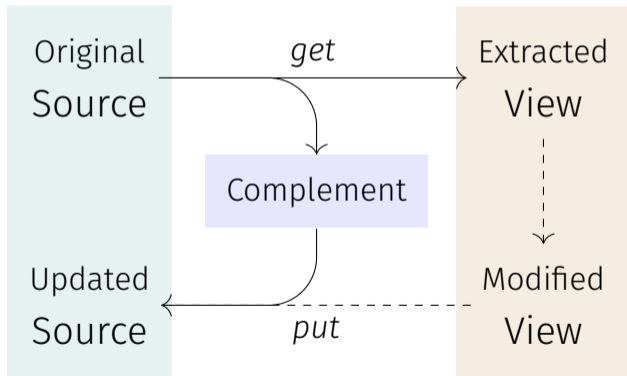
Bidirectional Transformation (Lens)

[Foster 2007; Hofmann et al. 2011]

Parsing/printing aside, we already know how to synchronise data.

Source (S)

View (V)



Non-Injectivity

`get` is not injective.

Cannot `put` with view alone.

Complement

Save the lost information.

`put` with complement.

Biparsers: Generalisation of Lenses

Lenses are pure; they cannot handle parsing/printing.

For composable parsing/printing, we need the **P monad** (state + error).

```
record P (A : Set) : Set where
  field runP : String → Result Error (A × String)
```

Biparsers are **lenses** generalised with **monadic parsing and printing**:

```
record Biparser (C X Y : Set) : Set where
  field parse : X → P (Y × C)
  print : Y × C → P X
```

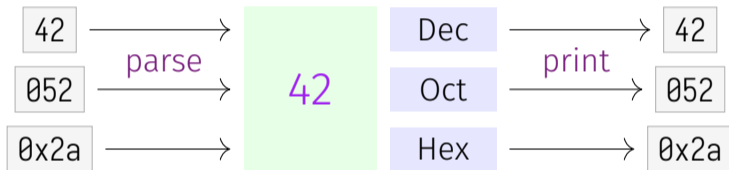
```
get : X → Y × C
put : Y × C → X
```

Existing definitions and properties for lenses also need to be generalised.

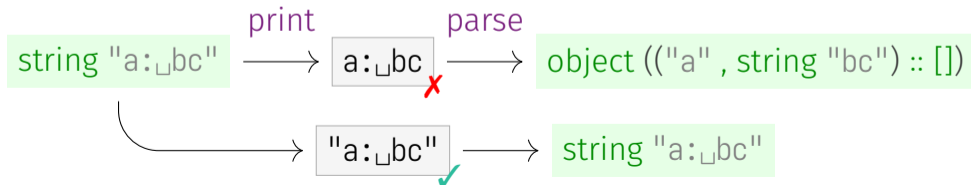
Generalised Round-Trip Properties

Generalising round-trip properties \Leftrightarrow addressing the two challenges.

- Challenge 1 (Exact-Printing): parsed data can be printed back

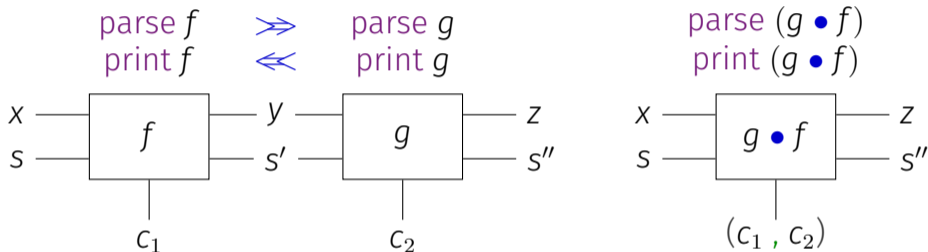


- Challenge 2 (Consistency): printed text can be parsed back



Generalised Composition

Composition uses Kleisli composition and accumulates the complement:



- Composition = Sequencing (p then q) + Transformation ($x \Rightarrow y \Rightarrow z$).
- Main result: **composition preserves round-trip properties.**

Compositional: combine small biparsers into larger biparsers.

Round-trip properties are automatically derived.

A Rich Combinator Language

Biparser Combinators

Classical parser combinator framework (e.g., Parsec) [Wadler 1995]:

- *item*: token consumption
- Sequencing: first p then q
- Alternation (Choice): either p or q Special Care for Round-Trip Properties
- Filtering: validate and transform Composition = Sequencing + Transform

By analogy, we define biparser combinators.

Main Challenges

- Take complements into consideration
- Make sure combinators preserve round-trip properties

Example: JSON Strings (Parser Combinators)

$\langle \text{String} \rangle ::= \text{"} \langle \text{StringChar} \rangle^* \text{"}$

Parser resembles the grammar

Sequencing

Alternation

Transformation

$\text{string} = \text{char } \text{'\"'} * \text{many stringChar} < * \text{char } \text{'\"'}$

$\langle \text{StringChar} \rangle ::= \text{NONCONTROL} \mid \backslash \langle \text{CEsc} \rangle \mid \backslash \langle \text{UniEsc} \rangle$

$\text{stringChar} = \text{nonControl} < | > (\text{char } \text{'\\'} * \text{cEsc}) < | > (\text{char } \text{'\\'} * \text{uniEsc})$

$\langle \text{CEsc} \rangle ::= \text{"} \mid \backslash \mid / \mid \text{b} \mid \text{f} \mid \text{n} \mid \text{r} \mid \text{t}$

$\begin{aligned} \text{cEsc} = & (\text{const } \text{'\"'} \langle \$ \rangle \text{char } \text{'\"'}) < | > (\text{const } \text{'\\'} \langle \$ \rangle \text{char } \text{'\\'}) \\ & < | > (\text{const } \text{'/'} \langle \$ \rangle \text{char } \text{'/'}) < | > (\text{const } \text{'\b'} \langle \$ \rangle \text{char } \text{'b'}) \\ & < | > (\text{const } \text{'\f'} \langle \$ \rangle \text{char } \text{'f'}) < | > (\text{const } \text{'\n'} \langle \$ \rangle \text{char } \text{'n'}) \\ & < | > (\text{const } \text{'\r'} \langle \$ \rangle \text{char } \text{'r'}) < | > (\text{const } \text{'\t'} \langle \$ \rangle \text{char } \text{'t'}) \end{aligned}$

$\langle \text{UniEsc} \rangle ::= \text{u} \text{ HEX HEX HEX HEX}$

$\text{uniEsc} = \text{char } \text{'u'} * (\text{mkUnicodeChar} \langle \$ \rangle \text{hex} < * \text{hex} < * \text{hex} < * \text{hex})$

Example: JSON Strings (Biparsers)

Biparser resembles the grammar

Biparser resembles the parser

Simple Development and Migration

```
⟨String⟩ ::= " ⟨StringChar⟩* "
```

```
string = char '"' #*>> many stringChar <<*> char '"'
```

```
⟨StringChar⟩ ::= NONCONTROL | \ ⟨CEsc⟩ | \ ⟨UniEsc⟩
```

```
stringChar = nonControl <+> (char '\\ ' #*>> cEsc) <+> (char '\\ ' #*>> uniEsc)
```

```
⟨CEsc⟩ ::= " | \ | / | b | f | n | r | t
```

```
cEsc = (const' '"' #• char '"') <+> (const' '\\ ' #• char '\\ ')  
<+> (const' '/' #• char '/') <+> (const' '\\b' #• char 'b')  
<+> (const' '\\f' #• char 'f') <+> (const' '\\n' #• char 'n')  
<+> (const' '\\r' #• char 'r') <+> (const' '\\t' #• char 't')
```

```
⟨UniEsc⟩ ::= u HEX HEX HEX HEX
```

```
uniEsc = char 'u' #*>> (mkUnicodeChar #• (hex <<*> hex <<*> hex <<*> hex))
```

Highlights of Our Solution

Compatibility between Lenses and Effects

Previous Work [Abou-Saleh et al., 2015]

Serious restrictions on how effects can be used in bidirectional programs.

We solve the problem for parser/printer effects:

Complement-Based Lenses

- Avoids the *get* in lens compositions.
- Makes composition preserve the round-trip properties.
- Exposes complements for reasoning; improves printing performance.

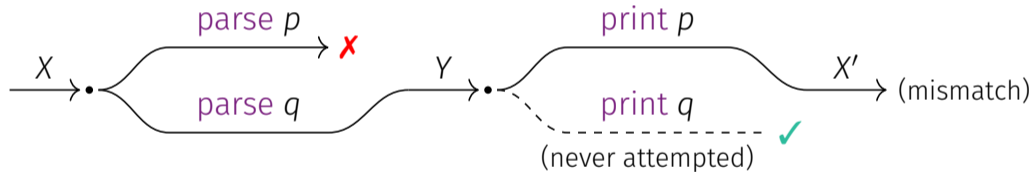
Effect Cancellation

- *print* *rewinds* the effects performed by *parse* (and vice versa).
- Fits our purpose perfectly and simplifies the theory.

Backtracking Choice and Consistency Hazard

Parser combinators rely on (shallow) backtracking for choices.

“ p or q ”: attempt p , and if p fails, attempt q .



Solution: different choice combinators with different pre-conditions.

More Notable Features

- Printing from scratch: optional complement.
- Context sensitivity: support for arbitrary user state.
- Improving output quality: view-complement alignment.
- Fine-grained complement manipulation.

Proof and Example

Agda Proof and Haskell Example

Mechanised Proof in Agda

- Around 1300 LoC (including comments and blank lines).
- All combinators with round-trip properties verified.

Runnable Examples in Haskell (JSON and YAML Subsets)

- Around 1000 LoC for combinators, 300 LoC for each example.
- All combinators with full Haddock documentation.

Both available as reusable packages.

Conclusion

Conclusion

- Biparsers with *exact-printing*: lens + monadic parsing/printing.
 - Formalised consistency: round-trip properties.
 - Compositional programming:
Composition preserves round-trip properties.
- Rich combinator language: biparser combinators.

Artifact available as reusable packages:

- Mechanised proof in Agda for the combinators.
- Runnable examples in Haskell for JSON and YAML.

More details in the paper!



Outline for Navigation

Outline

Problem: Data Synchronisation

Solution: Biparsers

A Rich Combinator Language

Highlights of Our Solution

Proof and Example

Conclusion

Details

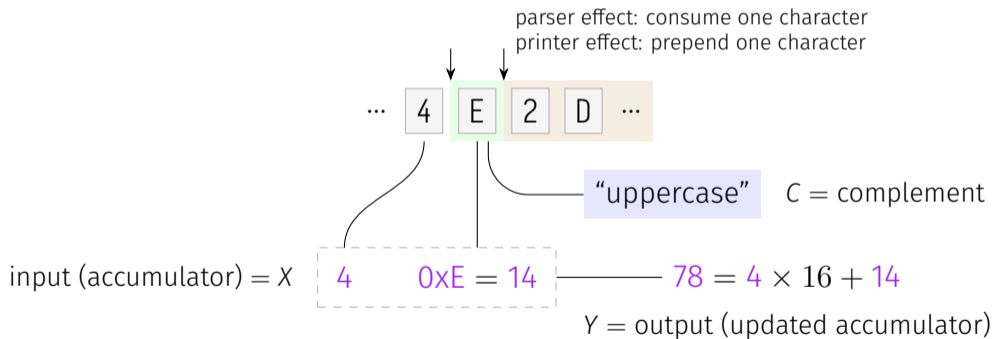
FAQ

1. What are X , Y , and C in practical biparsers?
2. Round-Trip Properties: What happens when the view is modified?
View-Complement Alignment: How does it work?
3. Can you elaborate on your choice combinators?
4. How do you design the complements?
Do you define a data type for the complement of every biparser?

Biparsers: Intuition

Two aspects of biparsers: parsing/printing + transformation.

Consider a biparser for hexadecimal numbers.



Round-Trip Properties, for Modified Text/Data

Edits do not void the guarantee of round-trip properties.

Power of Compositional Programming

- Biparsers are constructed compositionally.
- Component biparsers synchronise between **text segments** and **subtrees**.
- Round-trip properties apply to unchanged segments and subtrees.

Can we do even better?

- Before printing, perform a view-complement alignment.
- Generally applicable to any biparser.
- Various reusable strategies, orthogonal to parsing/printing.
- Only affects output quality, correctness always guaranteed.

Choice and the Consistency Conditions

Two Types of Backtracking Choices

- Biased choice: left branch preferred (the *regular* version).
- Remembered choice: the branch used last time is preferred.

Consistency Conditions

- Non-Intersection: disjoint branches (= unambiguous grammar).
- Consistency: same behaviour for intersection.
- Weak Consistency: consistency, under certain circumstances.

Complement Types

Complement types are not designed, but discovered.

- Follow the grammar and write the biparser.
- The suitable complement emerges as a result.

Fine-grained control over complements is also possible:

- Combinators dedicated to complement manipulation: [mapC](#) etc.
- Define complement data types for readability.